

Symbolic methods in statistics: elegance towards efficiency

E. Di Nardo

University of Basilicata, Italy

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Algebraic Statistics 2015



- ① What I mean by *symbolic methods*?
- ② Why symbolic methods?
- ③ The moment symbolic method
- ④ Applications to random matrices

How to convince people that in programming, simplicity and clarity - in short what mathematicians call elegance - are a crucial matter that decides between success and failure? (E. Dijkstra)

Outline

What I mean by *symbolic methods*?

Why symbolic methods?

The moment symbolic method

Applications to random matrices

A look into the past

An example

Throw away paper and pencil

The symbolic approach to combinatorial enumerations

The symbolic approach to combinatorial enumerations

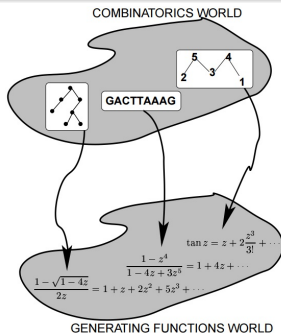
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The symbolic approach to combinatorial enumerations

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The symbolic approach to combinatorial enumerations

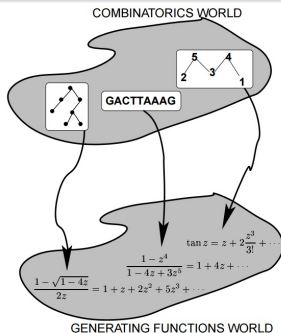


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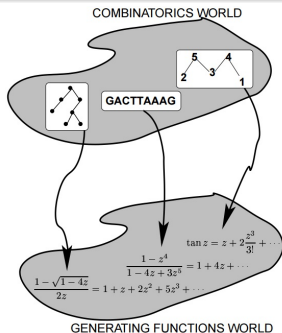
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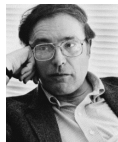


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- ▷ Flajolet, P. and Sedgewick, R. (2009) [Analytic Combinatorics](#), Cambridge Univ. Press
- ▷ Roman, S.M. and Rota, G.-C. (1978) [The umbral calculus](#), Adv.Math.

$$G(a_n; t) = \sum_{n \geq 0} a_n \frac{t^n}{n!} \text{ (exponential generating function)}$$

the “magic art” of lowering and raising exponents: $a_n \rightarrow a^n$



Gian-Carlo Rota

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The umbral calculus: an approach to combinatorial sequences using cumulants.

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Although the notation satisfied the most ardent advocate of spic-and-span rigor, the translation of “classical” umbral calculus into the newly found rigorous language made the method altogether unwieldy and unmanageable. Not only was the eerie feeling of witchcraft lost in the translation, but, after such a translation, the use of calculus to simplify computation and sharpen our intuition was lost by the wayside.

- ▷ Rota, G.-C. and Bryan, D.T. (1994) [The classical umbral calculus](#), *SIAM J. Math. Anal. Appl.*



- ▷ Bell, E.T. (1938) [The history of Blissards symbolic method with a sketch of its inventors life](#), *Amer. Math. Monthly*

What I mean by *symbolic methods*

A set of manipulation techniques aiming to perform algebraic calculations (possibly) through an algorithmic approach in order to find efficient mechanical processes to pass to a computer.

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Séminaire Lotharingien de Combinatoire

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What about computation cost?

- Efficiency is not so obvious.
- Sometimes a consequence of a different viewpoint.

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Why not a symbolic
cumulant calculus?



Very close to the moment
method for random matrices.



Commutative counterpart of
free probability.

U -statistics

An appropriate choice of language and notation can simplify and clarify many statistical calculations.

▷ McCullagh, P. (1987) **Tensor Methods in Statistics**. Chapman & Hall, London

U -statistics

$$U = \frac{1}{(n)_m} \sum \Phi(X_{j_1}, X_{j_2}, \dots, X_{j_m})$$

- symmetric polynomial in (X_1, X_2, \dots, X_n) i.i.d.r.v.'s;
- the sum ranges over the set of all permutations (j_1, j_2, \dots, j_m) .

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Augmented symmetric polynomials vs moments

If $\lambda = (1^{r_1}, 2^{r_2}, \dots, m^{r_m}) \vdash n$ is a partition of $r_1 + 2r_2 + \dots + mr_m = n$ of length $r_1 + r_2 + \dots + r_m = l(\lambda)$ and $E[X_i^j] = a_j$, $j = 1, 2, \dots, m$ then

$$E \left[\sum \underbrace{X_s X_t \dots}_{r_1} \underbrace{X_q^2 X_r^2 \dots}_{r_2} \dots \underbrace{X_u^m X_v^m \dots}_{r_m} \right] = (n)_{l(\lambda)} a_1^{r_1} a_2^{r_2} \dots a_m^{r_m}$$

▷ Stuart, A. and Ord, J.K. (1994) **Kendall's Advanced Theory of Statistics. Vol. 1: Distribution Theory** Edward Arnold, London (Section 12.5)

A key tool: the singleton umbra

A key tool: the singleton umbra

$$\sum \underbrace{X_s X_t \cdots}_{r_1} \underbrace{X_q^2 X_r^2 \cdots}_{r_2} \underbrace{X_u^m X_\nu^m \cdots}_{r_m} \cdots$$

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assume that \uparrow

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$$\prod_{j=1}^m (\chi_1 X_1^j + \chi_2 X_2^j + \cdots + \chi_n X_n^j)^{r_j}$$

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$$\prod_{j=1}^m (\chi_1 X_1^j + \chi_2 X_2^j + \cdots + \chi_n X_n^j)^{r_j}$$

having a structure very similar to

$$a_1^{r_1} a_2^{r_2} \cdots a_m^{r_m}$$

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How? $\triangleright \mathbb{E}[\chi_j^i] = \begin{cases} 1 & i = 0, 1 \\ 0 & \text{otherwise} \end{cases}$

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Are $\{\chi_i\}_{i=1}^n$ r.v.'s?

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 Singleton umbra

From vectors...



What calculations can be automated?

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 How can we automate them?

From vectors...

- What calculations can be automated?
- How can we automate them?
- What new concepts are required? (if any)

From vectors...

- ⚡→ What calculations can be automated?
- ⚡→ How can we automate them?
- ⚡→ What new concepts are required? (if any)

Computing

$$E \left[\left(\sum_{i \neq j}^n X_i^2 X_j \right) \left(\sum_{i=1}^n X_i^2 Y_i \right)^2 \right] \text{ with } (X_1, Y_1), \dots, (X_n, Y_n) \text{ separately i.i.d.r.v.'s}$$

▶ Push

From vectors...

- What calculations can be automated?
- How can we automate them?
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Computing

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► Push ... and then: (with $\mu_{i,j} = E[X^i Y^j]$)

$$2 \binom{n}{2} [2 \mu_{4,1} \mu_{3,1} + \mu_{5,2} \mu_{2,0} + \mu_{6,2} \mu_{1,0}] + 2 \binom{n}{3} \mu_{3,1} \mu_{2,1} \mu_{2,0} + \binom{n}{3} [2 \mu_{4,1} \mu_{2,1} \mu_{1,0} + \mu_{4,2} \mu_{2,0} \mu_{1,0}] + \binom{n}{4} \mu_{2,1}^2 \mu_{2,0} \mu_{1,0}$$

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- In a reasonable amount of time
- In a form suitable for any symbolic language

...up to random matrices

Computing

▶ Algorithm

$$E \left\{ \text{Tr} [W_p(n)H_1] \text{Tr} [W_p(n)H_2]^2 \right\} \quad \text{with } H_1, H_2 \in \mathbb{C}^{p \times p}$$

$$W_p(n) = \sum_{i=1}^n (\mathbf{X}_i - \mathbf{m}_i)^\dagger (\mathbf{X}_i - \mathbf{m}_i) \quad \text{and } \mathbf{X}_i \sim N(\mathbf{m}_i, \Sigma)$$

<Wishart random matrix>

...up to random matrices

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<Wishart random matrix>

▶ Push

... and then: $\left(\text{with } \Omega = \Sigma^{-1}M \text{ and } M = \sum_{i=1}^n \mathbf{m}_i^\dagger \mathbf{m}_i \right)$

$$\begin{aligned} & E \left\{ \text{Tr} [W_p(n)H_1] \text{Tr} [W_p(n)H_2]^2 \right\} = n \text{Tr} (H_2) \text{Tr} (\Omega H_1 H_2) - n \text{Tr} (H_2) \text{Tr} (\Omega H_2 H_1) \\ & + n \text{Tr} (H_2) \text{Tr} (\Omega H_1) \text{Tr} (\Omega H_2) - n \text{Tr} (\Omega H_2) \text{Tr} (H_1 H_2) - n^2 \text{Tr} (\Omega H_2) \text{Tr} (H_1) \text{Tr} (H_2) \\ & + 2 \text{Tr} (\Omega H_2) \text{Tr} (\Omega H_1 H_2) + 2 \text{Tr} (\Omega H_2) \text{Tr} (\Omega H_2 H_1) - \text{Tr} (\Omega H_1) (\text{Tr} (\Omega H_2))^2 \\ & - \text{Tr} (\Omega H_1 H_2^2) - \text{Tr} (\Omega H_2 H_1 H_2) + \text{Tr} (\Omega H_1) \text{Tr} (\Omega H_2^2) + 2 n^2 \text{Tr} (H_1 H_2) \text{Tr} (H_2) \\ & + n^2 / 2 \text{Tr} (H_1) \text{Tr} (H_2^2) + n^3 \text{Tr} (H_1) (\text{Tr} (H_2))^2 + n \text{Tr} (H_1 H_2^2) \\ & + n \text{Tr} (H_1) (\text{Tr} (\Omega H_2))^2 + n^2 \text{Tr} (\Omega H_1) (\text{Tr} (H_2))^2 - n / 2 \text{Tr} (\Omega H_1) \text{Tr} (H_2^2) \\ & - n \text{Tr} (H_1) \text{Tr} (\Omega H_2^2) \end{aligned}$$

Maple application center

The screenshot shows the Maple application center interface. At the top, there is a navigation bar with links for PRODUCTS, SOLUTIONS, PURCHASE, SUPPORT, RESOURCES, and COMPANY, along with a search bar. The main content area features a sidebar on the left with navigation options like Home, Editor's Choice Applications, and Application Search. The central content area displays a featured application titled "A new algorithm for computing moments of complex non-central Wishart distributions". The application details include authors (Dr. Giuseppe Guarino and Dr. Erika Di Nardo), publication date (February 25, 2013), and category (Mathematics: Abstract Algebra). There are also links for downloading the file, previewing the application, and contacting the author. A community rating section is visible on the right.

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Application Details

Authors: [Dr. Giuseppe Guarino](#)
[Dr. Erika Di Nardo](#)

Application Type: [Maple Document](#)

Publish Date: February 25, 2013

Created In: [Maple 12](#)

Language: English

Category: [Mathematics: Abstract Algebra](#)

Toolkit

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▷ Di Nardo, E. (2014) *On a symbolic representation of non-central Wishart random matrices with applications*. *Jour. Mult. Anal.*

In the literature

- ▷ Kendall, W.S. (1993) *Computer Algebra in probability and statistics*.
Statistica Neerlandica.

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- ▷ Rose, C. (2015) *MathStatca: a symbolic approach to computational mathematical statistics*. Version 2.7
- ▷ McCullagh, P. (1987) *Tensor methods in statistics*. Chapman and Hall.

In the literature

A steep learning curve but...

When symbolic methods are used properly, they can give us more insights to problems and the efficiency could be reached as by-product.

- ▷ Kendall, W.S. (1993) *Computer Algebra in probability and statistics*. Statistica Neerlandica.
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• mean $\rightarrow c_1$

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- mean $\rightarrow c_1$
- variance $\rightarrow c_2$
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- kurtosis $\rightarrow c_4/c_2^2$

k-statistics

Definition

The n -th k -statistic k_n is the unique symmetric unbiased estimator of the n -th cumulant c_n , i.e. $E[k_n] = c_n$.

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$$k_1 = \frac{S_1}{n}$$

$$k_2 = \frac{nS_2 - S_1^2}{(n)_2}$$

$$k_3 = \frac{2S_1^3 - 3nS_1S_2 + n^2S_3}{(n)_3}$$

$$k_4 = \frac{-6S_1^4 + 12nS_1^2S_2 - 3n(n-1)S_2^2 - 4n(n+1)S_1S_3 + n^2(n+1)S_4}{(n)_4}$$

$$S_r = \sum_{i=1}^n X_i^r$$

k -statistics

A nice formula: cumulants in terms of moments

If $\lambda = (1^{r_1}, 2^{r_2}, \dots, m^{r_m}) \vdash i \leq n$ then

$$c_i = i! \sum_{\lambda \vdash i} \frac{(-1)^{l(\lambda)-1} [l(\lambda) - 1]!}{r_1! r_2! \cdots r_m!} \frac{a_1^{r_1} a_2^{r_2} \cdots a_m^{r_m}}{(1!)^{r_1} (2!)^{r_2} \cdots (m!)^{r_m}}$$

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$$\frac{\sum \underbrace{X_s X_t \cdots X_q}_{r_1} \underbrace{X_q^2 X_r^2 \cdots X_u^m}_{r_2} \underbrace{X_u^m X_\nu^m \cdots}_{r_m} \cdots}{n(n-1) \cdots (n-l(\lambda)+1)}$$

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in terms of power sums

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▷ Too heavy from a computational point of view!

Computational results

(A&S) Andrews, D.F. and Stafford, J.E. (2000) *Symbolic computation for statistical inference*. Oxford University Press.

k -statistics	A&S
k_5	0.06
k_7	0.31
k_9	1.44
k_{11}	8.36
k_{14}	396.39
k_{16}	57982.4
k_{18}	—
k_{20}	—
k_{22}	—
k_{24}	—
k_{26}	—
k_{28}	—

PC Pentium(R)4, Intel(R)

CPU 2.08 Ghz

512MB Ram

Maple 10.0

Mathematica 4.2

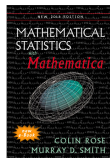
Times in seconds

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(A&S) Andrews, D.F. and Stafford, J.E. (2000) *Symbolic computation for statistical inference*. Oxford University Press.

(Symbolic) Di Nardo, E., Guarino, G. and Senato, D. (2008) *A unifying framework for *k*-statistics, polykays and their multivariate generalizations*. Bernoulli.

(MathStat)



<i>k</i> -statistics	A&S	MathStat	Symbolic
k_5	0.06	0.01	0.01
k_7	0.31	0.02	0.01
k_9	1.44	0.04	0.01
k_{11}	8.36	0.14	0.01
k_{14}	396.39	0.64	0.02
k_{16}	57982.4	2, 63	0.08
k_{18}	–	6.90	0.16
k_{20}	–	25.15	0.33
k_{22}	–	81.70	0.80
k_{24}	–	359.40	1.62
k_{26}	–	1581.05	2.51
k_{28}	–	6505.45	4.83

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A speedier way of computing

$$c_i = E[(C_{1,Z} + \cdots + C_{n,Z})^i]$$

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- $E[Z^m] = (-1)^{m-1} (m-1)! / (n)_m$ for $m = 0, 1, \dots, n$

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A speedier way of computing = a new formula and a new insight

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Statistics and Computing (2009)

“Every polynomial symmetric function can be expressed in terms of polykays.”

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$k_{r,\dots,s}$	AS Algorithms	MathStatica		Polyk-algorithm
$k_{3,2}$	0.06	0.02		0.02
$k_{4,4}$	0.67	0.06		0.06
$k_{5,3}$	0.69	0.08		0.07
$k_{7,5}$	34.23	0.79		0.70
$k_{7,7}$	435.67	2.52		2.43
$k_{9,9}$	-	27.41		23.32
$k_{10,8}$	-	30.24		25.06
$k_{4,4,4}$	34.17	0.64		0.77

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▷ Staude, B. and Rotter, S. (2010) *Cubic: cumulant based inference of higher-order correlations in massively parallel spike trains*. J. Comp. Neuroscience

Joint cumulants are zero for i.r.v.'s

Table: For AS Algorithms, missed computational times means “greater than 20 hours”. For MathStatistica, missed computational times means “procedures not available”

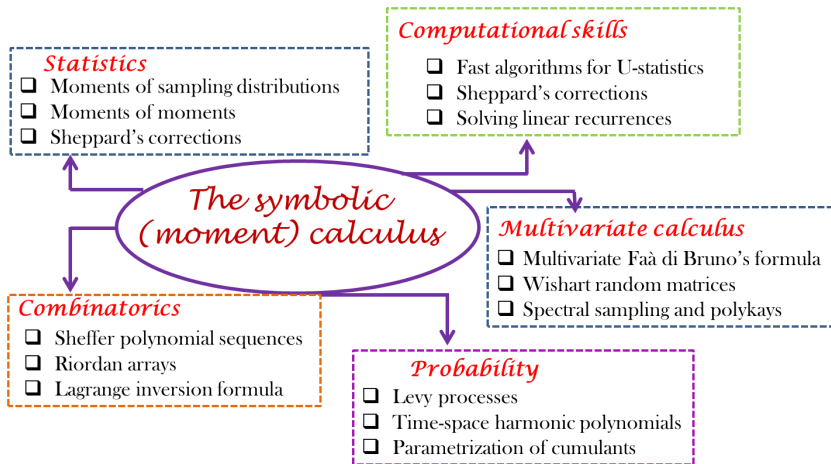
$k_{s_1 \dots s_r; l_1 \dots l_m}$	AS Algorithms	MathStatistica	
$k_{3 \ 2}$	0.25	0.03	
$k_{4 \ 4}$	28.36	0.16	
$k_{5 \ 5}$	259.16	0.55	
$k_{6 \ 5}$	959.67	1.01	
$k_{7 \ 7}$	-	8.49	
$k_{8 \ 7}$	-	14.92	
$k_{3 \ 3 \ 3}$	1180.03	0.88	
$k_{4 \ 4 \ 3}$	-	4.80	
$k_{4 \ 4 \ 4}$	-	13.53	
$k_{2 \ 1; \ 1 \ 1}$	0.20	-	
$k_{2 \ 2; \ 2 \ 1}$	6.30	-	
$k_{2 \ 2; \ 2 \ 2}$	33.75	-	
$k_{2 \ 2; \ 2 \ 1; \ 1 \ 1}$	126.19	-	
$k_{2 \ 2; \ 2 \ 1; \ 2 \ 1}$	398.42	-	
$k_{2 \ 2; \ 2 \ 2; \ 2 \ 1}$	1387.00	-	
$k_{2 \ 2; \ 2 \ 2; \ 2 \ 2}$	3787.41	-	

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$k_{4 \ 4 \ 3}$	-	4.80	0.94
$k_{4 \ 4 \ 4}$	-	13.53	2.30
$k_{2 \ 1; \ 1 \ 1}$	0.20	-	0.01
$k_{2 \ 2; \ 2 \ 1}$	6.30	-	0.08
$k_{2 \ 2; \ 2 \ 2}$	33.75	-	0.14
$k_{2 \ 2; \ 2 \ 1; \ 1 \ 1}$	126.19	-	0.28
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$k_{2 \ 2; \ 2 \ 2; \ 2 \ 1}$	1387.00	-	1.25
$k_{2 \ 2; \ 2 \ 2; \ 2 \ 2}$	3787.41	-	2.91

An overview on what we have done...



Symbolic moment calculus

a_i represented
by a symbol
 $\alpha \in \mathcal{A}$



$$\mathbb{E}[\alpha^i] = a_i$$

Symbolic combinatorics

$$a_i = |\{ \text{discrete structures} \}|$$

gen.func. $1 + \sum_{i \geq 1} a_i \frac{t^i}{i!}$



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Probability in terms of r.v.'s

Take an ordered commutative algebra over $\mathbb{C}[\mathcal{A}]$ and endow it with a positive linear functional \mathbb{E} :

- an element of $\mathcal{A} \Rightarrow$ a r.v.
- the linear functional $\mathbb{E} \Rightarrow$ the expectation of a r.v.
- the sequence $\{a_i\} \Rightarrow$ the moments of a r.v.

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Special Umbrae	Moments
Augmentation umbra	$\mathbb{E}[\varepsilon^i] = 0$
Unity umbra	$\mathbb{E}[u^i] = 1$
Singleton umbra	$\mathbb{E}[\chi^i] = \delta_{i,1}$
Bell umbra	$\mathbb{E}[\beta^i] = \mathfrak{B}_i$

Symbolic combinatorics

$$a_i = |\{\text{discrete structures}\}|$$

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Take an ordered commutative algebra over $\mathbb{C}[\mathcal{A}]$ and endows it with a positive linear functional \mathbb{E} :

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- the linear functional $\mathbb{E} \Rightarrow$ the expectation of a r.v.
- the sequence $\{a_i\} \Rightarrow$ the moments of a r.v.

Special Umbrae

Moments

Augmentation umbra

$\mathbb{E}[\varepsilon^i] = 0$

Unity umbra

$\mathbb{E}[u^i] = 1$

Singleton umbra

$\mathbb{E}[\chi^i] = \delta_{i,1}$

Bell umbra

$\mathbb{E}[\beta^i] = \mathfrak{B}_i$

▷ Not all r.v.'s can be represented by umbrae.

▷ Not all umbrae are r.v.'s.

The algebra of non-commutative r.v.'s

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 &\quad \text{Free cumulants - } \text{Nc}(i)
 \end{aligned}$$

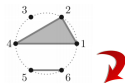
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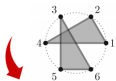


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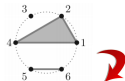
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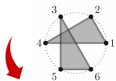
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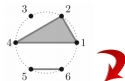
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Remark: a, b are free commutative r.v.'s iff at least one of them has vanishing variance.

The framework

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\mathbb{E} factorizes on different symbols

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The moment symbolic method

Applications to random matrices

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Again on efficiency: Wishart random matrices

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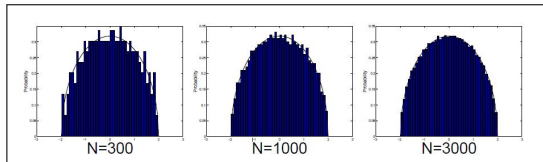
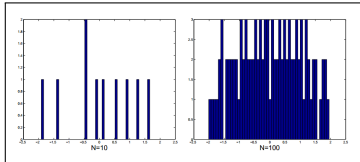
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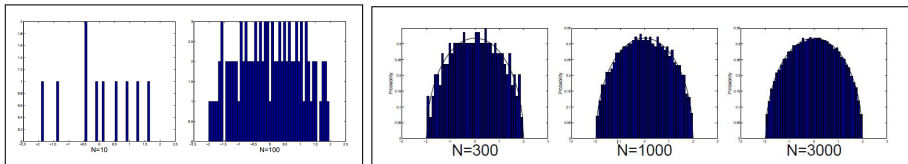
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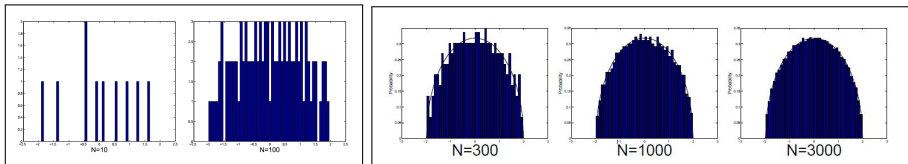
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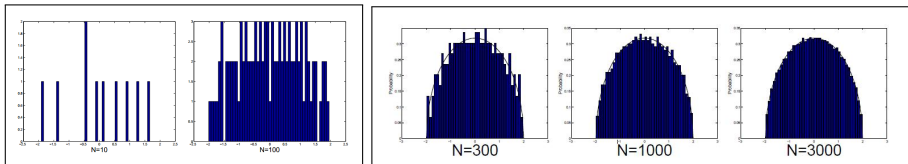
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- What about $c_i(A + B)$?
- How to define $c_i(A)$?

Non-asymptotic case

If A and B are asymptotically free, then the asymptotic spectrum of the sum can be obtained from the individual asymptotic spectra.



⚡ → As free probability only covers the asymptotic regime in which n is sent to infinity, there are some aspects of random matrix theory to which the tools of free probability are not sufficient by themselves to resolve.

How to preserve the framework of free probability?

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If in ~~$\lim_{n \rightarrow \infty} \frac{1}{n} E[\text{Tr}(A^i)] = \tau(A^i)$~~ the symbolic moment method can be resorted in order to compute $\{\tau(A^i)\}_{i \geq 1}$.

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Assume to symbolically represent the eigenvalues $\{\lambda_1, \dots, \lambda_n\}$ of A with $\{\mu_1, \dots, \mu_n\}$ umbral monomials so that

$$\tau(A^i) = \frac{1}{n} \mathbb{E}[\mu_1^i + \dots + \mu_n^i] \text{ power sum symmetric polynomials in } \{\mu_i\}$$

- ▷ Capitaine M., Casalis M. (2006)
Cumulants for random matrices as convolutions on the symmetric group.
Probab. Theory Relat. Fields.

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Cumulants of random matrices

If $\boldsymbol{\mu} = (\mu_1, \dots, \mu_n)$ represents the eigenvalues of A then $\mathbf{c}_{\boldsymbol{\mu}} = (c_{1,\boldsymbol{\mu}}, \dots, c_{n,\boldsymbol{\mu}})$

$$\text{Tr}(A) \Leftarrow \mu_1 + \dots + \mu_n \equiv n \cdot \beta \cdot (c_{1,\boldsymbol{\mu}} + \dots + c_{n,\boldsymbol{\mu}}) \Rightarrow \text{Tr}(\mathfrak{C}(A))$$

represents the n -tuple of cumulants of A .

▷ Capitaine M., Casalis M. (2006) *Cumulants for random matrices as convolutions on the symmetric group*. *Probab. Theory Relat. Fields*.

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$$\mathbf{m}[A(\sigma)] = \mathbb{E} \left\{ \prod_{c \in \mathcal{C}(\sigma)} \text{Tr} \left[A^{1(c)} \right] \right\} \quad \mathbf{c}[A(\sigma)] = \prod_{c \in \mathcal{C}(\sigma)} \frac{\mathbb{E} \{ [\text{Tr}(\mathfrak{C}(A))]^{1(c)} \}}{(1(c) - 1)!} \Rightarrow \text{polykeys}$$

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$$\alpha \equiv u \cdot \beta \cdot \kappa \quad \curvearrowright$$

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The non central Wishart distribution

- Let $\{\mathbf{X}_1, \dots, \mathbf{X}_n\}$ be random row vectors independently drawn from a p -variate complex normal distribution with zero mean and full rank covariance matrix Σ with eigenvalues $\{\theta_1, \dots, \theta_p\}$
- Let $\mathbf{m}_1, \dots, \mathbf{m}_n$ be complex row vectors of dimension p .



$$W_p(n, \Sigma, M) = \sum_{i=1}^n (\mathbf{X}_i - \mathbf{m}_i)^\dagger (\mathbf{X}_i - \mathbf{m}_i)$$

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Symbolic representation with $\widehat{W}(n) = W_p(n, \Sigma, 0)$

$$\text{Tr}[\widehat{W}(n)] \Leftarrow \text{central comp.} \\ + \underbrace{n \cdot (\theta_1 \bar{u}_1 + \dots + \theta_p \bar{u}_p)}$$

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$$\text{Cum}_i(\text{Tr}[W(n)]) = -i! \text{Tr}(M \Sigma^{i-1}) + n(i-1)! \text{Tr}(\Sigma^i)$$

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A different way to represent the central component

$$\mathrm{Tr}[\widehat{W}(n)] = \mathrm{Tr}[\mathbf{X}_1^\dagger \mathbf{X}_1 + \cdots + \mathbf{X}_n^\dagger \mathbf{X}_n]$$

with $\{\mathbf{X}_1^\dagger \mathbf{X}_1, \dots, \mathbf{X}_n^\dagger \mathbf{X}_n\}$ i.i.d. random matrices of order p .

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$$\mathrm{Tr}[\widehat{W}(n)] \equiv n.(\theta_1 \bar{u}_1 + \cdots + \theta_p \bar{u}_p) \equiv n.\beta.\delta$$

- ▷ $\{\bar{u}_1, \dots, \bar{u}_p\}$ uncorrelated umbrae similar to the **boolean unity** umbra \bar{u} whose moments are equal to the number of permutations of a set.

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As a summation of compound Poisson r.v.'s

$$\mathrm{Tr}[\widehat{W}(1)] = \mathrm{Tr} \left[\sum_{i=1}^n \mathbf{X}_i^\dagger \mathbf{X}_i \right] = Z_1 + \cdots + Z_{\mathrm{Po}(1)} \text{ with}$$

- $\{Z_i\}_{i=1}^n$ i.i.d. r.v.'s;
- $E[Z_i^k] = (k-1)! \mathrm{Tr}(\Sigma^k) = \mathrm{Cum}_k(\mathbf{X}_i^\dagger \mathbf{X}_i)$ for $k \in \mathbb{N}$

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↪ The sequence of moments of $\text{Tr}[\widehat{W}(n)]$ is of binomial type.

↪ The sequence of moments of $\text{Tr}[W(n)]$ is of Sheffer type.


Generalizing the computation of $m[W(n)]$ with multivariate notations

$$E \left\{ \text{Tr} [W(n)H_1]^{i_1} \cdots \text{Tr} [W(n)H_m]^{i_m} \right\}$$

[▶ Solution](#)

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$$E \left\{ \text{Tr} [W(n)H_1]^{i_1} \cdots \text{Tr} [W(n)H_m]^{i_m} \right\} = \mathbb{E} [(-1.\beta.\tilde{\eta} + n.\beta.\tilde{\rho})^i] \quad \text{▶ Solution}$$

 **Univariate case:** $E \left\{ \text{Tr} [W(n)]^k \right\} = \mathbb{E} [(-1.\beta.\alpha + n.\beta.\delta)^k]$

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Multivariate moments: *recepte ingredients*

For $\{g_i\}_{i \in \mathbb{N}_0^m} \in \mathbb{C}$ with $g_i = g_{i_1, i_2, \dots, i_m}$ and $g_0 = 1$, such that $\mathbb{E}[\nu^i] = g_i$

- $\nu = (\nu_1, \dots, \nu_m)$ m -tuple of umbral monomials (not necessarily uncorrelated)
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Univariate case: $E \left\{ \text{Tr} [W(n)]^k \right\} = \mathbb{E} [(-1.\beta.\alpha + n.\beta.\delta)^k]$

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▷ **Multinomial expansion:**
$$\sum_{\substack{t_1, t_2 \in \mathbb{N}_0^m \\ t_1 + t_2 = i}} \binom{i}{t_1, t_2} \mathbb{E} [(-1.\beta.\tilde{\eta})^{t_1}] \mathbb{E} [(n.\beta.\tilde{\rho})^{t_2}]$$

Generalizing the computation of $m[W(n)]$ with multivariate notations

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$$N_m[i] = \{ \text{necklaces of type } i \text{ on } [m] \} \begin{cases} N_3[(3, 0, 0)] & = \{111\} \\ N_3[(1, 2, 0)] & = \{122\} \\ N_3[(1, 1, 1)] & = \{123, 132\} \end{cases}$$

Generalizing the computation of $m[W(n)]$ with multivariate notations

$$E \left\{ \text{Tr} [W(\gamma)H_1]^{i_1} \cdots \text{Tr} [W(\gamma)H_m]^{i_m} \right\} = \mathbb{E} [(-1.\beta.\tilde{\eta} + \gamma.\beta.\tilde{\rho})^{\mathbf{i}}] \quad \text{▶ Solution}$$

Multivariate moments: *recipe ingredients*

For $\{g_i\}_{i \in \mathbb{N}_0^m} \in \mathbb{C}$ with $g_i = g_{i_1, i_2, \dots, i_m}$ and $g_0 = 1$, such that $\mathbb{E}[\nu^i] = g_i$

- $\nu = (\nu_1, \dots, \nu_m)$ m -tuple of umbral monomials (not necessarily uncorrelated)
- $\mathbf{i} \in \mathbb{N}_0^m$ multi-index.



Univariate case: $E \left\{ \text{Tr} [W(\gamma)]^k \right\} = \mathbb{E} [(-1.\beta.\alpha + \gamma.\beta.\delta)^k]$

▶ **Multinomial expansion:**
$$\sum_{\substack{t_1, t_2 \in \mathbb{N}_0^m \\ t_1 + t_2 = \mathbf{i}}} \binom{\mathbf{i}}{t_1, t_2} \mathbb{E} [(-1.\beta.\tilde{\eta})^{t_1}] \mathbb{E} [(\gamma.\beta.\tilde{\rho})^{t_2}]$$

$$N_m[\mathbf{i}] = \{ \text{necklaces of type } \mathbf{i} \text{ on } [m] \} \begin{cases} N_3[(3, 0, 0)] & = \{111\} \\ N_3[(1, 2, 0)] & = \{122\} \\ N_3[(1, 1, 1)] & = \{123, 132\} \end{cases}$$

Tricking: an example

$$\text{Cum}_{\mathbf{i}}(\text{Tr}[W(n) H_1], \dots, \text{Tr}[W(n) H_m]) = \mathbf{i}!(n\mathbb{E}[\rho^{\mathbf{i}}] - \mathbb{E}[\eta^{\mathbf{i}}])$$



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$$\sum_{\boldsymbol{\lambda} \vdash \mathbf{i}} \frac{\mathbb{E}[(\chi.n)^{l(\boldsymbol{\lambda})}]}{m(\boldsymbol{\lambda})} \prod_{\lambda_j} \mathbb{E}[\boldsymbol{\rho}^{\lambda_j}]^{r_j} \quad \text{(the central part)}$$

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Employ the results of $E \left\{ \text{Tr} [\widehat{W}(n) H_1] \text{Tr} [\widehat{W}(n) H_2]^2 \right\}$

$$2n^2 \text{Tr}(H_1 H_2) \text{Tr}(H_2) + n^3 \text{Tr}(H_1) (\text{Tr}(H_2))^2 + n \text{Tr}(H_1 H_2^2) + n^2 / 2 \text{Tr}(H_1) \text{Tr}(H_2^2)$$

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Replace n^i with c_i

Simple random sampling

Simple random sample

A sub-vector \mathbf{y} consisting of m components of $\mathbf{x} = (x_1, \dots, x_n) \in \mathbb{R}^n$, selected with equal probability $1/\binom{n}{m}$.

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$\sigma \in \mathfrak{S}_n$ a permutation

S the corresponding matrix

$$S_{ij} = \begin{cases} 1, & \text{if } \sigma(i) = j, \\ 0, & \text{otherwise.} \end{cases}$$

$$S = \begin{pmatrix} s_{1,1} & s_{1,2} & \dots & s_{1,n} \\ \vdots & \vdots & \dots & \vdots \\ s_{m,1} & s_{m,2} & \dots & s_{m,n} \\ s_{m+1,1} & s_{m+1,2} & \dots & s_{m+1,n} \\ \vdots & \vdots & \vdots & \vdots \\ s_{n,1} & s_{n,2} & \dots & s_{n,n} \end{pmatrix}$$

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↪ Properties:

$$\triangleright S_{[m \times n]} S_{[m \times n]}^T = I_m, \quad \triangleright S_{[m \times n]}^T S_{[m \times n]} \neq I_m.$$

Example

$$S = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix} \quad \text{corresponding to} \quad \sigma = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 4 & 3 & 1 & 2 \end{pmatrix} \in \mathfrak{S}_4$$

Example

$$S_2 = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix} \quad \text{corresponding to} \quad \sigma = \begin{pmatrix} 1 & 2 & 3 & 4 \\ 4 & 3 & 1 & 2 \end{pmatrix} \in \mathfrak{S}_4$$



A simple random sampling is:

$$\begin{pmatrix} x_4 & 0 \\ 0 & x_3 \end{pmatrix} = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix} \begin{pmatrix} x_1 & 0 & 0 & 0 \\ 0 & x_2 & 0 & 0 \\ 0 & 0 & x_3 & 0 \\ 0 & 0 & 0 & x_4 \end{pmatrix} \begin{pmatrix} 0 & 0 \\ 0 & 0 \\ 0 & 1 \\ 1 & 0 \end{pmatrix}$$

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The full matrix is:

$$\begin{pmatrix} x_4 & 0 & 0 & 0 \\ 0 & x_3 & 0 & 0 \\ 0 & 0 & x_1 & 0 \\ 0 & 0 & 0 & x_2 \end{pmatrix} = \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \end{pmatrix} \begin{pmatrix} x_1 & 0 & 0 & 0 \\ 0 & x_2 & 0 & 0 \\ 0 & 0 & x_3 & 0 \\ 0 & 0 & 0 & x_4 \end{pmatrix} \begin{pmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 \end{pmatrix}$$

Outline

What I mean by *symbolic methods*?

Why symbolic methods?

The moment symbolic method

Applications to random matrices

What free probability can do for statistician?

Again on efficiency: Wishart random matrices

Spectral random sampling

Conclusions

Spectral sample

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Let H be a random unitary matrix uniformly distributed with respect to the Haar measure on the group \mathcal{U}_n of $n \times n$ unitary matrices.

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$$\mathbf{x} \rightarrow \boxed{H_{n-m}} \xrightarrow{Y} \text{red arrow} Y = H_{[m \times n]} \text{diag}(\mathbf{x}) H_{[m \times n]}^\dagger$$

Spectral sample of size m

The eigenvalues (real r.v.'s)

$$\mathbf{y} = (y_1, \dots, y_m) \text{ of } Y$$

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⚡ **Generalization:** replace $\text{diag}(\mathbf{x})$ with a Hermitian random matrix X

⚡ **Meaning:** a restriction operation $X \mapsto Y$ extracting a partial information from X

A second meaning

A random Hermitian matrix A of order n is said to be *freely randomized* if its distribution is invariant under unitary conjugation, i.e. $A \sim GAG^\dagger$ for each unitary G .

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- if $m = n$, the subsample \mathbf{y} is a random permutation of \mathbf{x} .
- if $m < n$, the elements of \mathbf{y} do not occur among the components of \mathbf{x} .

Statistics for spectral sampling?



Which spectral properties are preserved on the average by freely randomized matrix restriction? **Example:** the eigenvalue average is preserved.

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▷ Tukey, J. (1956) *Keeping moment-like sample computations simple*. Ann.Math.Stat.

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$$E [T_m(\mathbf{y})|\mathbf{x}] = T_n(\mathbf{x}) \text{ for each } m \leq n \text{ and } \mathbf{y} \text{ drawn from } \mathbf{x}$$

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$$E [T_m(\mathbf{y})|\mathbf{x}] = T_n(\mathbf{x}) \text{ for each } m \leq n \text{ and } \mathbf{y} \text{ drawn from } \mathbf{x}$$

- symmetric functions for $m = n$;
- not a single function in isolation, but a list of functions;
- common interpretation independent of the sample size

Statistics for spectral sampling?

 Which spectral properties are preserved on the average by freely randomized matrix restriction? **Example:** the eigenvalue average is preserved.

▷ Tukey, J. (1956) *Keeping moment-like sample computations simple*. Ann.Math.Stat.


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Spectral natural statistics

If \mathbf{y} spectral sample and $\lambda \vdash i$, then $E[\kappa_\lambda(\mathbf{y})] = \prod_{j=1}^{l(\lambda)} E[(c_{1,\mathbf{y}} + \dots + c_{m,\mathbf{y}})^{\lambda_j}]$

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Matricial polykeys

Main theorem

Matricial polykeys are the symmetric functions $\mathfrak{K}_\lambda(\mathbf{y})$ such that

$$\mathbb{E}[\mathfrak{K}_\lambda(\mathbf{y})](\sigma) = \text{const} \times \mathbb{E} \left\{ \left[\mu(I_m)^{(-1)} \star \mu(Y) \right] (\sigma) \right\}, \quad i \leq m$$

- $\sigma \in \mathfrak{S}_m$, a permutation with $|C(\sigma)|$ disjoint cycles;
- $(f \star g)(\sigma) = \sum_{\rho \omega = \sigma} f(\rho) g(\omega)$ convolution on \mathfrak{S}_m ;
- $\mu(Y)(\sigma) = \prod_{c \in C(\sigma)} \text{Tr} \left(Y^{I(c)} \right)$ and $\mu(I_m)(\sigma) = m^{|C(\sigma)|}$;
- $f^{(-1)} \star f = f \star f^{(-1)} = \delta$ (indicator function)

The computation of $\mu(I_m)^{(-1)}$ requires to solve a system of m equations in m indeterminates. A different way: the so-called **Weingarten function** on \mathfrak{S}_m (**Open problem**).

$$\mathfrak{K}_{(1)} = \frac{S_1}{n} \quad \mathfrak{K}_{(2)} = \frac{nS_2 - S_1^2}{n(n^2 - 1)}$$

$$\mathfrak{K}_{(1^2)} = \frac{nS_1^2 - S_2}{n(n^2 - 1)}$$

$$\mathfrak{K}_{(3)} = 2 \frac{2S_1^3 - 3nS_1S_2 + n^2S_3}{n(n^2 - 1)(n^2 - 4)}$$

$$\mathfrak{K}_{(1,2)} = \frac{-2nS_3 + (n^2 + 2)S_1S_2 - nS_1^3}{n(n^2 - 1)(n^2 - 4)}$$

$$\mathfrak{K}_{(1^3)} = \frac{S_1^3(n^2 - 2) - 3nS_1S_2 + 4S_3}{n(n^2 - 1)(n^2 - 4)}$$

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Connection with k -statistics

$$\mathfrak{K}_{(1)} = \frac{S_1}{n} = k_{(1)} \quad \mathfrak{K}_{(2)} = \frac{nS_2 - S_1^2}{n(n^2 - 1)} = \frac{k_{(2)}}{(n+1)}$$

$$\mathfrak{K}_{(1^2)} = \frac{nS_1^2 - S_2}{n(n^2 - 1)} = \frac{k_{(1^2)}}{(n+1)}$$

$$\mathfrak{K}_{(3)} = 2 \frac{2S_1^3 - 3nS_1S_2 + n^2S_3}{n(n^2 - 1)(n^2 - 4)} = \frac{2k_{(3)}}{(n+1)(n+2)}$$

$$\mathfrak{K}_{(1,2)} = \frac{-2nS_3 + (n^2 + 2)S_1S_2 - nS_1^3}{n(n^2 - 1)(n^2 - 4)} = \frac{2k_{(1,2)} - nk_{(1)}k_{(2)}}{(n+1)(n+2)}$$

$$\mathfrak{K}_{(1^3)} = \frac{S_1^3(n^2 - 2) - 3nS_1S_2 + 4S_3}{n(n^2 - 1)(n^2 - 4)} = \frac{2k_{(1^3)} - 3k_{(1)}k_{(2)} + n(n+3)(k_{(1)})^3}{(n+1)(n+2)}.$$

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Properties

- $\mathfrak{R}_\lambda(\mathbf{y})$ are called **matricial polykays**, **unbiased estimators** of products of cumulants
- $\mathfrak{R}_\lambda(\mathbf{y})$ are **natural statistics** (the proof is strictly connected with the spectral sampling);
- The condition $i \leq m$ parallels the analogous condition for Fisher's k -statistics.
- $\mathbb{E}[\mathfrak{R}_\lambda(\mathbf{y})]$ tends towards the **product of free cumulants** when $m \rightarrow \infty$ as Fisher's polykays tends towards the product of classical cumulants.
- $\mathfrak{R}_\lambda(\mathbf{y})$ can be expressed as linear combination of generalized k -statistics with coefficients independent of n .

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Generalized spectral polykeys

Generalized k -statistics are the sample version of the generalized cumulants.

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Generalized cumulants

$$c_{r,s,t} = \text{cov}(X^r, X^s X^t) \quad c_{r,s,t,u} = \text{cov}(X^r X^s, X^t X^u)$$


 **application:** in asymptotic approximations of distributions

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Generalized k -statistics are the sample version of the generalized cumulants.

- ↪ the generalized k -statistics are linearly independent;
- ↪ every polynomial symmetric function can be expressed uniquely as a linear combination of generalized k -statistics;
- ↪ any polynomial symmetric function whose expectation is independent of n can be expressed as linear combination of generalized k -statistics with coefficients independent of n .

A different choice of foundations can lead to a different way of thinking about the subject, and thus to ask a different set of questions and to discover a different set of proofs and solutions. Thus it is often of value to understand multiple foundational perspectives at once, to get a truly stereoscopic view of the subject.

From Terence Tao's blog

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([Topics on random matrices, Terence Tao](#)) - on line

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Thanks for your attention!

Cumulants:

theory, computation and applications ↪

Work in progress: Wiley

